

# MATH 2200 - EXAM 1 SPRING 2017

## SOLUTION

Thursday 9 February 2017  
Instructor: Tom Cuchta

### Instructions:

- Show all work, clearly and in order, if you want to get full credit. If you claim something is true **you must show work backing up your claim**. I reserve the right to take off points if I cannot see how you arrived at your answer (even if your final answer is correct).
- Justify your answers whenever possible to ensure full credit.
- Circle or otherwise indicate your final answers.
- Please keep your written answers brief; be clear and to the point.
- Good luck!

Tautological implications and equivalences

**Commutative Law of  $\vee$ :**  $P \vee Q \leftrightarrow Q \vee P$

**Law of Equivalence for Implication and Disjunction:**  $(\neg P \vee Q) \leftrightarrow (P \rightarrow Q)$

**Law of Contraposition:**  $(P \rightarrow Q) \leftrightarrow (\neg Q \rightarrow \neg P)$

**Modus tollendo ponens:**  $[(\neg P) \wedge (P \vee Q)] \rightarrow Q$

**Modus tollendo tollens:**  $[(\neg Q) \wedge (P \rightarrow Q)] \rightarrow \neg P$

**Law of Detachment:**  $[P \wedge (P \rightarrow Q)] \rightarrow Q$

1. (9 points) Translate the following sentences into symbolic notation, with letters standing only for atomic sentences. Be sure to identify which sentences correspond to which letters you use.

**This problem resembles #3(h) in HW1.**

- (a) (3 points) “The sum is positive or the sum is zero.”

*Solution:* Let  $P$  be the sentence “The sum is positive.” and let  $Q$  be the sentence “The sum is zero”. The given sentence, written symbolically, is  $P \vee Q$ .

- (b) (3 points) “If  $f$  is midpoint-convex and  $f$  is continuous, then  $f$  is convex.”

*Solution:* Let  $P$  be the sentence “ $f$  is midpoint-convex.” and let  $Q$  be the sentence “ $f$  is convex”. Then the given sentence, written symbolically, is  $P \rightarrow Q$ .

- (c) (3 points) “The graph  $G$  is a tree if and only if  $G$  is connected and  $G$  does not contain a cycle.”

*Solution:* Let  $P$  be the sentence “The graph  $G$  is a tree.”, let  $Q$  be the sentence “ $G$  does not contain a cycle.”, and let  $R$  be the sentence “ $G$  contains a cycle”. The given sentence, written symbolically, is  $P \leftrightarrow (Q \wedge \neg R)$ .

(note: if you chose  $R$  to instead be “ $G$  does not contain a cycle”, then you would have gotten  $P \leftrightarrow (Q \wedge R)$ , which is ok)

2. (16 points) If  $P$  and  $Q$  are true and  $R$  is false, then are the following sentences true or false?

**This problem resembles #4(a), #5(d), #6(e), and #6(i) in HW1.**

- (a) (4 points)  $P \vee R$

*Solution:* True (because  $P$  is true)

- (b) (4 points)  $\neg(\neg R)$

*Solution:* False (because  $\neg R$  is true, which is the case because  $R$  is false)

- (c) (4 points)  $R \rightarrow \neg P$

*Solution:* True (because “False  $\rightarrow$  anything” is always true)

- (d) (4 points)  $R \leftrightarrow [P \rightarrow (Q \wedge R)]$

*Solution:* True ( $Q \wedge R$  is false, meaning  $P \rightarrow (Q \wedge R)$  is false since  $P$  is true. This means both  $R$  and  $P \rightarrow (Q \wedge R)$  are both false, making the biconditional true, since they match.)

3. (12 points) Use a truth table to determine whether the given sentence is a tautology or not.

**This problem resembles #10(c) and #10(i) in HW2.**

- (a) (6 points)  $[P \wedge (\neg Q)] \vee Q$

*Solution:*

$P$	$Q$	$\neg Q$	$P \wedge (\neg Q)$	$[P \wedge (\neg Q)] \vee Q$
$T$	$T$	$F$	$F$	$T$
$T$	$F$	$T$	$T$	$T$
$F$	$T$	$F$	$F$	$T$
$F$	$F$	$T$	$F$	$F$

From the final column, we see that this is **not** a tautology.

- (b) (6 points)  $[\neg P \wedge (P \vee Q)] \rightarrow Q$

*Solution:*

$P$	$Q$	$\neg P$	$P \vee Q$	$[\neg P \wedge (P \vee Q)]$	$[\neg P \wedge (P \vee Q)] \rightarrow Q$
$T$	$T$	$F$	$T$	$F$	$T$
$T$	$F$	$F$	$T$	$F$	$T$
$F$	$T$	$T$	$T$	$T$	$T$
$F$	$F$	$T$	$F$	$F$	$T$

From the final column, we that this **is** a tautology.

4. (12 points) **This problem resembles #16(c) in HW2.**

(a) (6 points) Does the sentence  $P \wedge (P \rightarrow Q)$  tautologically imply the sentence  $Q$ ? If so, prove it using a truth table. If not, explain why not.

*Solution:* Yes it does. To see it, we must show that  $[P \wedge (P \rightarrow Q)] \rightarrow Q$  is a tautology.

$P$	$Q$	$P \rightarrow Q$	$P \wedge (P \rightarrow Q)$	$[(P \rightarrow Q)] \rightarrow Q$
$T$	$T$	$T$	$T$	$T$
$T$	$F$	$F$	$F$	$T$
$F$	$T$	$T$	$F$	$T$
$F$	$F$	$T$	$F$	$T$

(b) (6 points) Does the sentence  $\neg P \wedge (P \rightarrow Q)$  tautologically imply the sentence  $P$ ? If so, prove it using a truth table. If not, explain why not.

*Solution:* It does not. We show it via truth table:

$P$	$Q$	$\neg P$	$P \rightarrow Q$	$\neg P \wedge (P \rightarrow Q)$	$[\neg P \wedge (P \rightarrow Q)] \rightarrow P$
$T$	$T$	$F$	$T$	$F$	$T$
$T$	$F$	$F$	$F$	$F$	$T$
$F$	$T$	$T$	$T$	$T$	$F$
$F$	$F$	$T$	$T$	$T$	$F$

Since the last table shows that the implication  $[\neg P \wedge (P \rightarrow Q)] \rightarrow P$  is **not** a tautology, we say that  $\neg P \wedge (P \rightarrow Q)$  does **not** tautologically imply  $P$ .

5. (14 points) **This problem resembles #17(b) and #17(d) in HW2.**

(a) (7 points) Is the sentence  $Q \vee \neg P$  tautologically equivalent to  $Q \vee P$ ? If so, prove it using a truth table. If not, explain why not.

*Solution:* It is not. We demonstrate via truth table:

$P$	$Q$	$\neg P$	$Q \vee \neg P$	$Q \vee P$	$(Q \vee \neg P) \rightarrow (Q \vee P)$	$(Q \vee P) \rightarrow (Q \vee \neg P)$
$T$	$T$	$F$	$T$	$T$	$T$	$T$
$T$	$F$	$F$	$F$	$T$	$T$	$F$
$F$	$T$	$T$	$T$	$T$	$T$	$T$
$F$	$F$	$T$	$T$	$F$	$F$	$T$

Since the last two columns do not exactly match, we see that they are not tautologically equivalent. *note: we could have also made another column for  $(Q \vee P) \leftrightarrow (Q \vee \neg P)$  which would show that that sentence is not a tautology — this also shows the two sentences are not tautologically equivalent.*

(b) (7 points) Is the sentence  $P$  tautologically equivalent to the sentence  $\neg(\neg P)$ ? If so, prove it using a truth table. If not, explain why not.

6. (9 points) Make a counterexample to show the following rule of inference is invalid: “from  $\neg Q$  and  $P \rightarrow \neg Q$ , we may conclude  $P$ .”

**This problem resembles #1 in HW3.**

*Solution:* We must pick truth values for  $P$  and  $Q$  that make the premises true and the conclusion false. So we let  $Q$  be false (so that the premise  $\neg Q$  is true) and we will let  $P$  be false (so that the conclusion will be false). Note that doing this makes the other premise  $P \rightarrow \neg Q$  to also be false.

This means we have found a counterexample because our premises  $\neg Q$  and  $P \rightarrow \neg Q$  are true while simultaneously the conclusion  $P$  is false (based on how we chose the truth values for  $P$  and  $Q$ ).

7. (12 points) There **three** errors in the following formal deduction. Find them and explain why they are errors:

**This problem resembles #2 in HW3.**

{1}	(1) $P$	Premise
{2}	(2) $P \rightarrow \neg Q$	Premise
{3}	(3) $R \rightarrow Q$	Premise
{3}	(4) $\neg Q \rightarrow \neg R$	3 T (Law of Contraposition)
{1, 2}	(5) $\neg Q$	1, 2 T (Modus tollendo ponens)
{1, 2, 3, 4, 5}	(7) $\neg R$	4, 5 T (Law of Detachment)

*Solution:*

**First error:** Line (5) claims that it follows from lines 1 and 2 using modus tollendo ponens. This is not true, it actually follows from the Law of Detachment.

**Second error:** The left columns of the last line of the deduction claims that this line depends on the premises {1, 2, 3, 4, 5}, but lines 4 and 5 are not premises, and so they should not be listed here.

**Third error:** The last line of the deduction is called (7), but it should be (6).

8. (16 points) Construct a formal deduction for the following arguments. Do so by applying the stated rule(s).

**This problem resembles #4 and #5 in HW3.**

- (a) (8 points) “If  $n$  is odd, then  $i^n$  is complex. In actuality,  $i^n$  is not complex. Therefore  $n$  is not odd.”

**Rules:** Use **Modus tollendo tollens**.

*Solution:* Let  $P$  be the sentence “ $n$  is odd”, let  $Q$  be the sentence “ $i^n$  is complex”. We may write a formal deduction for this argument as follows:

{1}	(1) $P \rightarrow Q$	Premise
{2}	(2) $\neg Q$	Premise
{1, 2}	(3) $\neg P$	1, 2 T Modus tollendo tollens

- (b) (8 points) “If the axiom of choice holds, then nonmeasurable sets exist. If nonmeasurable sets exist, then the Banach-Tarski paradox is true. I will quit mathematics or the Banach-Tarski paradox is not true. The axiom of choice holds. Therefore I will quit mathematics.”

**Rules:** Use the **Law of Detachment** to conclude “the Banach-Tarski paradox is true” and then use **Commutative Law of  $\vee$**  to conclude “the Banach-Tarski paradox is not true or I will quit mathematics.” From here, use one more law (which?) to conclude “If the Banach-Tarski paradox is true, then I will quit mathematics.” After that, complete the deduction using the **Law of Detachment**.

*Solution:* Let

$P$  = “The axiom of choice holds.”

$Q$  = “Nonmeasurable sets exist.”

$R$  = “The Banach-Tarski paradox is true.”

$S$  = “I will quit mathematics.”

Construct the formal deduction as follows:

{1}	(1) $P \rightarrow Q$	Premise
{2}	(2) $Q \rightarrow R$	Premise
{3}	(3) $S \vee (\neg R)$	Premise
{4}	(4) $P$	Premise
{1, 4}	(5) $Q$	1, 4 T (Law of Detachment)
{1, 2, 4}	(6) $R$	2, 5 T (Law of Detachment)
{1, 2, 3, 4}	(7) $(\neg R) \vee S$	3 T (Commutative Law of $\vee$ )
{1, 2, 3, 4}	(8) $R \rightarrow S$	7 T (Law of Equivalence for Implication and Disjunction)
{1, 2, 3, 4}	(9) $S$	6, 8 T (Law of Detachment)