

A theory T is a set of sentences of first order logic (called axioms of the theory). Any sentence provable from axioms is called a theorem of T .

So far in this class, we have proved sentences in the "empty theory" ~ theory w/ no axioms.
 (e.g. any time you proved $\vdash S$)

First Order Arithmetic

A theory with some useful predicates:

\rightarrow predicates: $+$, \cdot , $S(x)$ S means "successor"
name: 0
 We will NOT have "1", "2", "3"
 Instead we have $0, S(0), S(S(0)), \dots$
names in theory

Side note: $a+b$ not $+(a,b)$

Axioms of First Order Arithmetic

(2)

① $\forall x \neg (0 = Sx)$

② $\forall x \forall y (Sx = Sy \rightarrow x = y)$

← encodes
" $x+1 = y+1 \rightarrow x=y$ "

③ $\forall y (y = 0 \vee \exists x (Sx = y))$

④ $\forall x (x + 0 = x)$

⑤ $\forall x \forall y (x + Sy = S(x+y))$

← encodes
 $x + (y+1) = (x+y) + 1$

⑥ $\forall x (x \cdot 0 = 0)$

⑦ $\forall x \forall y (x \cdot Sy = (x \cdot y) + x)$

← encodes
 $x \cdot (y+1) = x \cdot y + x$

⑧ $\forall x \forall y (x + y = y + x)$

⑨ $\forall x \forall y (x \cdot y = y \cdot x)$

↑
Standard interpretation

domain: $\{0, 1, 2, 3, \dots\}$

$+$: addition

\cdot : multiplication

Let's prove FOA \vdash $S0+S0=SS0$

(3)

1st order arithm

assume these premises & don't write them

- 1 $\forall y (S0 + Sy = S(S0 + y))$
- 2 $S0 + S0 = S(S0 + 0)$
- 3 $S0 + 0 = S0$
- 4 $S0 + S0 = SS0 \quad =E$

$x = S0$

$\forall E \text{ Ax } 5$

$\forall E \text{ 1}$

$y = 0$

$\forall E \text{ Ax } 4$

$=E \text{ 3, 2} \quad x = S0$

Prove

4

$$\text{FOA} \vdash \text{SSO} \cdot \text{SSO} = \text{SSSSO}$$

$x = \text{SSO}$

1 $\forall y (\text{SSO} \cdot S_y = (\text{SSO} \cdot y) + \text{SSO})$ $\forall E \text{ Ax } 7$

2 $\text{SSO} \cdot \text{SSO} = (\text{SSO} \cdot \text{SO}) + \text{SSO}$ $\forall E 1$
 $y = \text{SO}$

3 $\text{SSO} \cdot \text{SO} = (\text{SSO} \cdot 0) + \text{SSO}$ $\forall E 1$
 $y = 0$

4 $\text{SSO} \cdot 0 = 0$

$\forall E \text{ Ax } 6$

5 $\text{SSO} \cdot \text{SO} = 0 + \text{SSO}$

$x = \text{SSO}$
 $= E 4, 3$

6 $\text{SSO} + 0 = \text{SSO}$

$\forall E \text{ Ax } 4$

7 $\forall y (\text{SSO} + y = y + \text{SSO})$

$\forall E \text{ Ax } 8$

8 $\text{SSO} + 0 = 0 + \text{SSO}$

$\forall E 7$
 $y = 0$

9 $0 + \text{SSO} = \text{SSO}$

$= E 8, 6$

10 $\text{SSO} \cdot \text{SO} = \text{SSO}$

$= E 9, 5$

11 $\text{SSO} \cdot \text{SSO} = \text{SSO} + \text{SSO}$

$= E 10, 2$

12 $\forall y (\text{SSO} + S_y = S(\text{SSO} + y))$ $\forall E \text{ Ax } 5$
 $x = \text{SSO}$

13 $\text{SSO} + \text{SSO} = S(\text{SSO} + \text{SO})$ $\forall E 12$
 $y = \text{SO}$

14 $\text{SSO} + \text{SO} = S(\text{SSO} + 0)$ $\forall E 12$
 $y = 0$

15 $\text{SSO} + \text{SO} = \text{SSSSO}$

$= E 6, 14$

16 $\text{SSO} + \text{SSO} = \text{SSSSO}$

$= E 15, 13$

17 $\text{SSO} \cdot \text{SSO} = \text{SSSSO}$

$= E 16, 11$

$$\text{FOA} \vdash \forall x (x \cdot (x+SO) = x \cdot x + x) \quad x(x+1) = x^2 + x$$

(5)

$$1 \quad \forall y (a \cdot Sy = (a \cdot y) + a)$$

$$\forall E Ax 7 \quad x=a$$

$$2 \quad a \cdot Sa = (a \cdot a) + a$$

$$\forall E I \quad y=a$$

$$3 \quad \forall y (a + Sy = S(a + y))$$

$$\forall E Ax 5 \quad x=a$$

$$4 \quad a + SO = S(a + 0)$$

$$\forall E 3 \quad y=0$$

$$5 \quad a + 0 = a$$

$$\forall E Ax 4 \quad x=a$$

$$6 \quad a + SO = Sa$$

$$= E 5, 4$$

$$7 \quad a \cdot (a + SO) = a \cdot a + a$$

$$= E 6, 2$$

$$8 \quad \forall x (x \cdot (x + SO) = x \cdot x + x)$$

$$\forall I 7$$

Sets

Mathematical notion of a collection of objects.

Sets are "unordered collections"

Have 2 fundamental properties:

① two sets have same members iff they are the same set

② given any predicate P , the set of x in universe such that $P(x)$ is true should be a set

$\pi \in \mathbb{R}$

Axioms of (naive) set theory ~ predicate \in
"element of"

① Extensionality

$$\forall z (z \in x \leftrightarrow z \in y) \leftrightarrow x = y$$

② (Unrestricted) Comprehension (Axiom schema)

For any predicate P :

$$\exists z \forall x (x \in z \leftrightarrow P(x))$$